

CSL model checking

Quantitative Logics

15 May 2013

David N. Jansen

Recapitulation: CTMC

- A continuous-time Markov chain consists of:
 - S finite set of states
(often $S = \{1, 2, \dots, n\}$)
 - $\mathbf{R}: S \times S \rightarrow \mathbb{R}_{\geq 0}$ transition **rate** matrix
 - $\pi_0: S \rightarrow [0,1]$ initial state distribution
(sometimes)
 - $L: S \rightarrow AP$ labelling with atomic propositions

Recapitulation: CSL

- state formulas ϕ, ψ
 - a atomic proposition
 - $\neg\phi$ negation
 - $\phi \vee \psi$ disjunction
 - $\mathbf{P}_{\leq p}(\Pi), \mathbf{P}_{\geq p}(\Pi)$ probabilistic operator
 - $\mathbf{S}_{\leq p}(\phi), \mathbf{S}_{\geq p}(\phi)$ steady-state operator
- path formulas Π
 - $X^I \phi$ next state
 - $\phi U^I \psi$ until

with time bound:
an interval $I \subseteq \mathbb{R}_{\geq 0}$

CSL Model checking

- Assume given a CTMC (S, \mathbf{R}, L) and a formula ϕ
- Find, for each state formula ψ that is in ϕ , the set of states that satisfies it, $Sat(\psi)$
- Find, for each path formula Π that is in ϕ , the set of paths that satisfies it, $Sat(\Pi)$

Simple formulas

- Atomic proposition:

$$Sat(a) = \{ s \in S \mid a \in L(s) \}$$

- Negation:

$$Sat(\neg\phi) = S \setminus Sat(\phi)$$

- Disjunction:

$$Sat(\phi \vee \psi) = Sat(\phi) \cup Sat(\psi)$$

Next formulas

- $Sat(\mathbf{P}_{\leq p}(X^I \phi)) = \{ s \in S \mid Prob_s(Sat(X^I \phi)) \leq p \}$
- Do not actually calculate this satisfaction set!
- This is similar to PCTL next formula & probability to take the transition in interval I .

Until formulas

- $Sat(\mathbf{P}_{\leq p}(\phi U^I \psi))$
 $= \{ s \in S \mid Prob_s(Sat(\phi U^I \psi)) \leq p \}$
- There is not a fixed number of transitions in this interval.
- Use Fox–Glynn sum, combined with time-bounded until model checking

- we are going to make this slide together

Steady-state formulas

- $Sat(\mathbf{S}_{\leq p}(\phi)) =$
 $\{ s \in S \mid \pi_s \leq p \text{ for the equation } \pi Q = 0 \}$