

# Reductieregels voor Semantische Tableaus

Beweren en Bewijzen

28 maart 2006

$\neg_L :$ $\begin{array}{c} \Phi, \neg\alpha \circ \Psi \\   \\ \Phi \circ \alpha, \Psi \end{array}$	$\neg_R :$ $\begin{array}{c} \Phi \circ \neg\alpha, \Psi \\   \\ \Phi, \alpha \circ \Psi \end{array}$
$\wedge_L :$ $\begin{array}{c} \Phi, \alpha \wedge \beta \circ \Psi \\   \\ \Phi, \alpha, \beta \circ \Psi \end{array}$	$\wedge_R :$ $\begin{array}{c} \Phi \circ \alpha \wedge \beta, \Psi \\ / \quad \backslash \\ \Phi \circ \alpha, \Psi \quad \Phi \circ \beta, \Psi \end{array}$
$\vee_L :$ $\begin{array}{c} \Phi, \alpha \vee \beta \circ \Psi \\ / \quad \backslash \\ \Phi, \alpha \circ \Psi \quad \Phi, \beta \circ \Psi \end{array}$	$\vee_R :$ $\begin{array}{c} \Phi \circ \alpha \vee \beta, \Psi \\   \\ \Phi \circ \alpha, \beta, \Psi \end{array}$
$\rightarrow_L :$ $\begin{array}{c} \Phi, \alpha \rightarrow \beta \circ \Psi \\ / \quad \backslash \\ \Phi, \beta \circ \Psi \quad \Phi \circ \alpha, \Psi \end{array}$	$\rightarrow_R :$ $\begin{array}{c} \Phi \circ \alpha \rightarrow \beta, \Psi \\   \\ \Phi, \alpha \circ \beta, \Psi \end{array}$
$\leftrightarrow_L :$ $\begin{array}{c} \Phi, \alpha \leftrightarrow \beta \circ \Psi \\ / \quad \backslash \\ \Phi, \alpha, \beta \circ \Psi \quad \Phi \circ \alpha, \beta, \Psi \end{array}$	$\leftrightarrow_R :$ $\begin{array}{c} \Phi \circ \alpha \leftrightarrow \beta, \Psi \\ / \quad \backslash \\ \Phi, \alpha \circ \beta, \Psi \quad \Phi, \beta \circ \alpha, \Psi \end{array}$